

ONE  $\omega$ -INCONSISTENT FORMALIZATION OF SET THEORY  
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Let  $T$  be a formal system containing  $ZF$ , whose language include explicitly the arithmetic signature and numerical variables run over  $\omega$ . Moreover, assume that there are  $T$ -formulas of the form  $[\psi(\bar{x}), \Delta(Q, \bar{x}, a), a]$ , where  $a$  is a numerical variable,  $\psi$  is a formula of  $ZF$ -signature (enriched by numerical variables),  $\Delta$  a formula of the same signature with a predicate variable  $Q$ . Other  $T$ -formulas are constructed from  $ZF$ -formulas and formulas described above using propositional connectives and quantifies over individual variables. We consider these formulas as meaningless combinations of symbols.

We assume that  $T$ -axiomatics contains the following axiom schemes:

$$[\psi(\bar{x}), \Delta(Q, \bar{x}, a), a]_0^a \leftrightarrow \psi(\bar{x});$$

$$[\psi(\bar{x}), \Delta(Q, \bar{x}, a), a]_{b+1}^a \leftrightarrow (\Delta_{[\psi(\bar{x}), \Delta(Q, \bar{x}, a), a]_b^a}^{Q, a}).$$

Other  $T$ -axioms are  $ZF$ -axioms  $ZF$ -axiom schemes extended (as well as logical schemes) to arbitrary  $T$ -formulas. It is not hard to check that  $T$  is inconsistent wrt  $ZF + (\textit{existence of strongly inaccessible cardinals})$ .

**Theorem 0.1** *The system  $T$  is  $\omega$ -inconsistent.*

This fact can be established at metamathematical level. This means that for a suitable  $T$ -formula  $\varphi(a)$  one can find two metamathematical objects: a formal  $T$ -proof of the sentence  $\exists a \varphi(a)$  and a primitive recursive function  $G(a)$ , which construct for a given numeral  $n$  a Goedel number of a  $T$ -proof of  $\neg \varphi(n)$ , and this property of the function  $G$  can be proved in the recursive arithmetic.

From this fact follows, in particular, that the existence of strongly inaccessible cardinals is refutable in  $ZF$ .